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APPLICATION NO.	FILING DATE	FIRST NAMED INVENTOR	ATTORNEY DOCKET NO.	CONFIRMATION NO.
09/774,906	01/30/2001	Shahrokh Sadjadi	50325-0511	5989
29989	7590	04/21/2005	EXAMINER	
HICKMAN PALERMO TRUONG & BECKER, LLP 2055 GATEWAY PLACE SUITE 550 SAN JOSE, CA 95110			BOUTAH, ALINA A	
			ART UNIT	PAPER NUMBER
			2143	
				DATE MAILED: 04/21/2005

Please find below and/or attached an Office communication concerning this application or proceeding.

Office Action Summary	Application No.	Applicant(s)	
	09/774,906	SADJADI, SHAHROKH	

Examiner	Art Unit	
Alina N Boutah	2143	

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --
Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If the period for reply specified above is less than thirty (30) days, a reply within the statutory minimum of thirty (30) days will be considered timely.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

- 1) Responsive to communication(s) filed on 12 October 2004.
 2a) This action is FINAL. 2b) This action is non-final.
 3) Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

Disposition of Claims

- 4) Claim(s) 1-44 is/are pending in the application.
 4a) Of the above claim(s) 1-7, 31, 36 and 41 is/are withdrawn from consideration.
 5) Claim(s) _____ is/are allowed.
 6) Claim(s) 8-30, 32-35, 37-40 and 42-44 is/are rejected.
 7) Claim(s) _____ is/are objected to.
 8) Claim(s) _____ are subject to restriction and/or election requirement.

Application Papers

- 9) The specification is objected to by the Examiner.
 10) The drawing(s) filed on 30 January 2001 is/are: a) accepted or b) objected to by the Examiner.
 Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).
 Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
 11) The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. § 119

- 12) Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
 a) All b) Some * c) None of:
 1. Certified copies of the priority documents have been received.
 2. Certified copies of the priority documents have been received in Application No. _____.
 3. Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).

* See the attached detailed Office action for a list of the certified copies not received.

Attachment(s)

- 1) Notice of References Cited (PTO-892)
 2) Notice of Draftsperson's Patent Drawing Review (PTO-948)
 3) Information Disclosure Statement(s) (PTO-1449 or PTO/SB/08)
 Paper No(s)/Mail Date _____
- 4) Interview Summary (PTO-413)
 Paper No(s)/Mail Date. _____
- 5) Notice of Informal Patent Application (PTO-152)
- 6) Other: _____

DETAILED ACTION

Election/Restrictions

Applicant's election without traverse of Group B (claims 8-30, 32-35, 37-40 and 42-44) in the reply filed on October 12, 2004 is acknowledged.

Claim Rejections - 35 USC § 103

The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negated by the manner in which the invention was made.

Claims 8-30, 32-35, 37-40 and 42-44 are rejected under 35 U.S.C. 103(a) as being unpatentable over USPN 6,253,236 issued to Troxel et al. (hereinafter referred to as Troxel) in view of USPN 6,493,804 issued to Soltis et al. (hereinafter referred to as Soltis).

Regarding claim 8, Troxel teaches a method of controlling use by concurrent users of a distributed resource on a network, wherein use of the resource is limited to a specified maximum number of concurrent users, the method comprising the computer-implemented steps of: providing a lock manager process executing on a host (abstract; col. 1, lines 15-36); associating a user identification for each user with one host (figures 4A-D); and

responding to a request for the resource associated with a first user having a first user identification associated with a first host by requesting a lock from a first local lock manager process executing on the first host (abstract; figure 4D; col. 1, lines 15-36).

However, Troxel does not explicitly teach providing a distributed lock manager process comprising a plurality of local lock manager processes executing on a corresponding plurality of hosts; associating a user identification for each user with one host of the plurality of hosts; and responding to a request for the resource associated with a first user having a first user identification associated with a first host of the plurality of hosts by requesting a lock from a first local lock manager process executing on the first host.

Soltis teaches providing a distributed lock manager process comprising a plurality of local lock manager processes executing on a corresponding plurality of hosts (figure 4: clients 105A-N);

associating a user identification for each user with one host of the plurality of hosts (figure 4: clients 105A-N);

responding to a request for the resource associated with a first user having a first user identification associated with a first host of the plurality of hosts by requesting a lock from a first local lock manager process executing on the first host (col. 3, lines 41-46).

At the time the invention was made, one of ordinary skill in the art would have been motivated to combine the teachings of Troxel with the teaching of Soltis in order to provide decentralized control of shared data, therefore maintaining data consistency (col. 3, lines 41-45).

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Regarding claim 9, Soltis teaches a method as recited in Claim 8, wherein during said step of associating a user identification with one host, the first user is associated with the first host based on information indicating that the first user signs onto the network most frequently at a terminal node of the network that uses fewer intervening network devices to pass messages to the first host than to any other host of the plurality of hosts (col. 9, lines 42-60).

Regarding claim 10, Soltis teaches a method as recited in Claim 8, wherein during said step of associating a user identification with one host, the first user is associated with the first host based on information indicating that the first user signs onto the network most frequently at a terminal node of the network geographically closer to the first host than to any other of the plurality of hosts (col. 9, lines 42-60).

Regarding claim 11, Troxel teaches a method as recited in Claim 8, wherein each local lock manager process maintains a lock data structure associated with the resource and the lock data structure includes a local resource maximum field (abstract); and further comprising the steps of generating and storing a value in the local resource maximum field maintained by each local lock manager process such that a summation over all the local lock manager process of the value in the local resource maximum field does not exceed the maximum number of concurrent users (col. 5, lines 20-41).

Regarding claim 12, Troxel teaches a method as recited in Claim 11, wherein a first value in the local resource maximum field maintained by the first local lock manager process is based on a number of users associated with the first host (col. 5, lines 20-41).

Regarding claim 13, Soltis teaches a method as recited in Claim 8, wherein a copy of the distributed resource resides on a host of the plurality of hosts (col. 9, lines 29-41).

Regarding claim 14, Soltis teaches a method as recited in Claim 8, wherein a copy of the distributed resource resides on a computing device that is not among the plurality of hosts (col. 9, lines 29-41).

Regarding claim 15, Troxel teaches a method as recited in Claim 11, further comprising: determining whether a number of outstanding locks granted by the first local lock manager process is less than a particular value stored in the local resource maximum field maintained by the first local lock manager process; and if the number of outstanding locks is less than the particular value, then generating and returning a lock object (col. 5, lines 20-41).

Regarding claim 16, Troxel teaches a method of controlling concurrent users of a distributed resource on a network, the resource limited to a maximum number of concurrent users, the method comprising the computer-implemented steps of:

receiving a request for the distributed resource from a client process for a user having a user identification (figure 4D);

determining a home location associated with the user identification (abstract; col. 1, lines 56-66);

sending a request for a lock object for the distributed resource to a first local lock manager process of the distributed lock manager process, the request including the home location (col. 1, lines 36-41);

receiving the lock object for the distributed resource from a lock manager process executing on the host, if a number of outstanding locks granted by the local lock manager process is less than a value of a local resource maximum determined for the second local lock manager process (col. 5, lines 20-41); and

providing access to the resource to the first client only in response to receiving the lock object (abstract; figure 4D; col. 5, lines 20-41).

However, Troxel fails to teach the home location indicating a unique host among a plurality of hosts that execute a corresponding plurality of local lock manager processes of a distributed lock manager process.

Soltis teaches the home location indicating a unique host among a plurality of hosts that execute a corresponding plurality of local lock manager processes of a distributed lock manager process (figure 4: clients 105A-N).

At the time the invention was made, one of ordinary skill in the art would have been motivated to combine the teachings of Troxel with the teaching of Soltis in order to provide decentralized control of shared data, therefore maintaining data consistency (col. 3, lines 41-45).

Regarding claim 17, Soltis teaches a method as recited in Claim 16, wherein a first host on which the first local lock manager process executes is different from the unique host on which the second local lock manager process executes (figure 4).

Regarding claim 18, Soltis teaches a method as recited in Claim 16, wherein a first host on which the first local lock manager process is executing is in a local network with a computing device on which the client process is executing (figure 4).

Regarding claim 19, Soltis teaches a method as recited in Claim 16, further comprising the steps of determining whether a request from the client process to set the retry time period has been received, and if so, automatically repeating the steps of determining, sending, receiving and providing (abstract).

Claim 20 is similar to claim 16 therefore is also rejected under the same rationale.

Regarding claim 21, Troxel teaches a method as recited in Claim 20, further comprising the steps of: when the second lock manager process is not associated with the particular home location, determining whether a first number of outstanding locks granted by the first local lock manager process for the distributed resource is less than a first local maximum number of users no greater than the maximum number of concurrent users, and if so, then incrementing the first number of outstanding locks granted and providing the lock object to the resource server (col. 5, lines 20-41).

Regarding claim 22, Troxel teaches a method as recited in Claim 20, wherein each local lock manager process of the distributed lock manager process stores a corresponding local maximum number of users for the distributed resource in a lock data structure associated with the local lock manager process, and an aggregate of local maximum numbers in lock data structures associated with all lock manager processes in the distributed lock manager process for the distributed resource is no greater than the maximum number of concurrent users to which the distributed resource is limited (col. 5, lines 20-41).

Regarding claim 23, Troxel teaches a method as recited in Claim 20, further comprising the steps of: receiving a request for a lock object from a third local lock manager process of the distributed lock manager process in response to a request from a second resource server; determining whether a first number of outstanding locks granted by the first look manager process for the distributed resource is less than a first local maximum number of users no greater than the maximum number of concurrent users; and if it is determined the first number of outstanding locks granted is less than the first local maximum number, then incrementing the first number of outstanding locks granted, and providing the lock object for the second resource server (col. 5, lines 20-41).

Regarding claim 24, Troxel teaches a method as recited in Claim 20, further comprising the step of providing the lock object to the resource server in response to receiving the lock object from the second lock manager process (col. 5, lines 20-41).

Regarding claim 25, Troxel teaches a method of distributing a resource on a network, the resource limited to a maximum number of concurrent users, the method comprising the steps of:

providing a distributed lock manager executing on a corresponding host (abstract; col. 1, lines 15-36); and

generating a value for a local resource maximum number of users stored on the host such that a summation of the value for the local resource maximum yields an aggregate value that does not exceed the maximum number of concurrent users (abstract; figure 4D; col. 1, lines 15-36).

However, Troxel does not teach determining whether to increase a first value in a first resource maximum stored on a first host; and if it is determined to increase the first resource minimum, then decreasing by a particular amount a second value in a second resource maximum stored on a second host of the plurality of hosts, and increasing by the particular amount the first value in the first resource maximum stored on the first host, wherein each local lock manager process is configured to grant a lock for the resource if the number of outstanding locks granted by the local lock manager process is less than a value of the local resource maximum stored on the corresponding host.

Soltis teaches determining whether to increase a first value in a first resource maximum stored on a first host; and if it is determined to increase the first resource minimum, then decreasing by a particular amount a second value in a second resource maximum stored on a second host of the plurality of hosts, and increasing by the particular amount the first value in the first resource maximum stored on the first host, wherein each local lock manager process is configured to grant a lock for the resource if the number of outstanding locks granted by the

local lock manager process is less than a value of the local resource maximum stored on the corresponding host (figure 4: clients 105A-N).

At the time the invention was made, one of ordinary skill in the art would have been motivated to combine the teachings of Troxel with the teaching of Soltis in order to provide decentralized control of shared data, therefore maintaining data consistency (col. 3, lines 41-45).

Regarding claim 26, Soltis teaches A method as recited in Claim 25, said step of determining whether to increase the first value further comprising determining the time of day and the geographic location served by the first host (col. 9, lines 42-60).

Regarding claim 27, Soltis teaches a method as recited in Claim 25, said step of determining whether to increase the first value further comprising determining whether a difference between the number of outstanding locks granted and the first value is less than a first predetermined threshold (col. 9, lines 42-60).

Regarding claim 28, Soltis teaches a method as recited in Claim 25, further comprising, if it is determined to increase the first value, then determining whether to decrease the value based on a whether a difference between the number of outstanding locks granted and the second value is greater than a second predetermined threshold (col. 9, lines 42-60).

Regarding claim 29, Soltis teaches a method as recited in Claim 28, wherein the particular amount is based on the second predetermined threshold (col. 9, lines 42-60).

Regarding claim 30, Soltis teaches a method as recited in Claim 29, wherein the particular amount is no greater than the second predetermined threshold (col. 9, lines 42-60).

Claim 32 is similar to claim 8 therefore is rejected under the same rationale.

Claim 33 is similar to claim 16 therefore is also rejected under the same rationale.

Claims 34, 35, 37-39, 40 and 42-44 are similar to claim 25 therefore are also rejected under the same rationale.

Conclusion

Any inquiry concerning this communication or earlier communications from the examiner should be directed to Alina N. Boutah whose telephone number is 571-272-3908. The examiner can normally be reached on Monday-Friday (9:00 am - 5:00 pm).

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, David A. Wiley can be reached on 571-272-3923. The fax phone number for the organization where this application or proceeding is assigned is 703-872-9306.



DAVID WILEY
SUPERVISORY PATENT EXAMINER
TECHNOLOGY CENTER 2100

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